

Extended Spiking Neural P Systems Generating Strings and Vectors of Non-Negative Integers

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Extended Spiking Neural P Systems ESNP Systems

$$\Pi = (m, S, r)$$

- m **number of cells**
(neurons) uniquely identified by a number between 1 and m
- S **initial configuration**
initial value (of spikes) in each neuron
- R **finite set of rules** $(i, E, k, d; r)$

ESNP Systems: Rules

Rules: $(i, E/k, d; r)$ where

- $i \in [1..m]$ specifies cell i this rule is assigned to
- $E \in \text{REG}(N)$ checking set
- $k \geq 0$ "number of spikes" consumed by this rule
- $d \geq 0$ delay (omitted when being 0)
- r set of productions of the form (l, w, t)

where

- $l \in [1..m]$ specifies the target cell
- $w \geq 0$ weight of energy sent along the axon from neuron i to neuron l ,
- t delay along the axon (omitted when being 0)

ESNP Systems - Configuration

- for each **neuron**, the **actual number of spikes** in the neuron is specified;
- in each **neuron** i , an "**activated rule**" $(i, E/k, d'; r)$ may wait to be executed where d' is the remaining time until the neuron spikes;
- in each **axon** to a neuron l , we may find **pending packages** of the form (l, w, t') where t' is the remaining time until w spikes have to be added to neuron l provided it is not closed for input at the time this package arrives.

ESNP Systems – Transition 1

- for each neuron i

first check whether there is an "activated rule" $(i, E/k, d'; r)$ waiting to be executed;

if $d'=0$, then neuron i "spikes", i.e.,
for every production (l, w, t) in r the corresponding package (l, w, t) is put on the axon from neuron i to neuron l , and after that, this "activated rule" $(i, E/k, 0; r)$ is eliminated;

ESNP Systems – Transition 2

- for each neuron l ,
now consider all packages (l, w, t') on axons leading to neuron l ;
provided the neuron is not closed, all weights w in such packages where $t'=0$ are summed up and added to the corresponding number of spikes in neuron l ;

in any case, the packages with $t'=0$ are eliminated from the axons, whereas for all packages with $t'>0$, t' is decremented by one;

ESNP Systems – Transition 3

- for each neuron i ,
check again whether there is an "activated rule" $(i, E/k, d'; r)$ (with $d' > 0$) or not;
 - NO: any rule $(i, E/k, d; r)$ from R for which the current number of spikes in the neuron is in E can be applied (and then put a copy of this rule as "activated rule" for this neuron into the description of the current configuration);
 - YES: replace d' by $d'-1$ and keep $(i, E/k, d'-1; r)$ as the "activated rule" in neuron i in the description of the configuration for the next step of the computation.

ESNP Systems – SNP Systems

In the original model, in the productions (l, w, t) of a rule $(i, E/k, d; \{ (l, w, t) \})$, only $w=1$ (for **spiking rules**) or $w=0$ (for **forgetting rules**) as well as $t=0$ was allowed (and for forgetting rules, the checking set E had to be finite and disjoint from all other sets E in neuron i). Moreover, **reflexive axons**, i.e., leading from neuron i to neuron i , were not allowed.

Yet the most important extension is that **different rules for neuron i may affect different axons** leaving from it whereas in the original model the structure of the axons (called synapses there) was fixed.

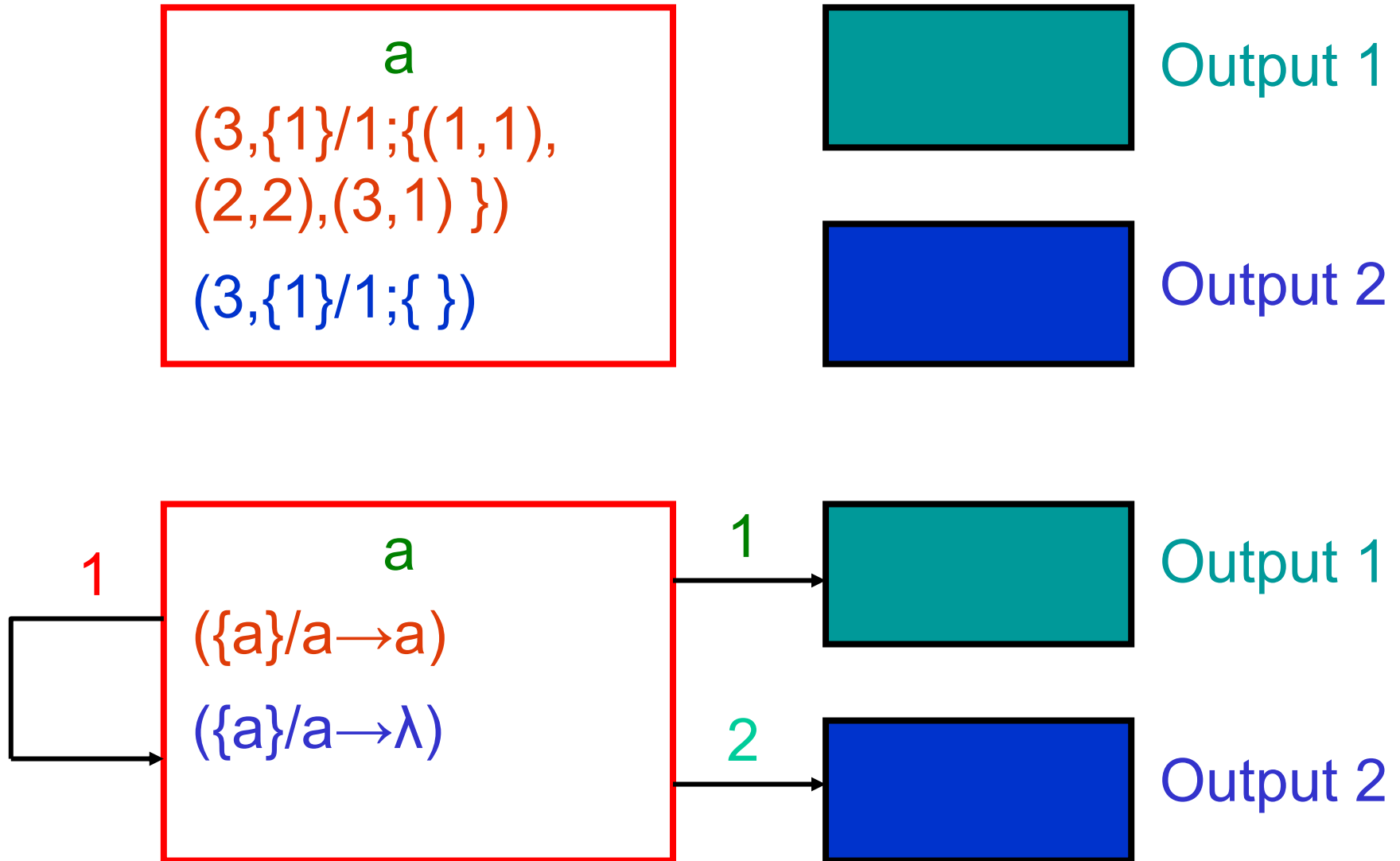
A First Example

$\{(m, 2m) \mid m \geq 0\}$ can easily be generated by an ESNP system *in real time* with only one *actor* neuron and two output neurons.

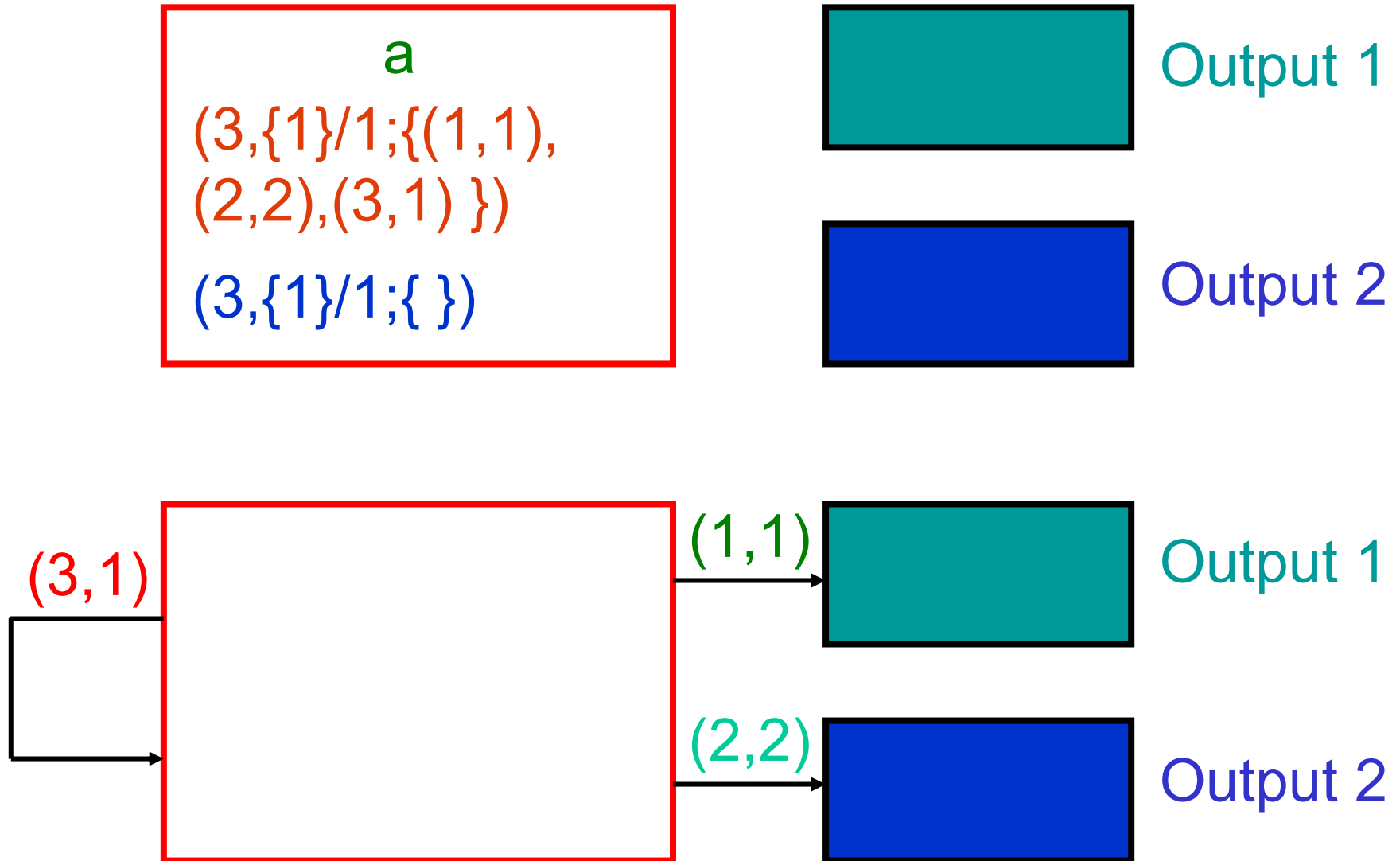
Other than in SNP systems

- at different times different numbers of spikes can be sent to various neurons.
- the output is taken as the number of spikes present in the output neurons at the end of a halting computation (and not as the time between the first two spikes in each output neuron).

A First Example – ESNP System



A First Example – Computation



A First Example – Computation

a
(3,{1}/1;{(1,1),
(2,2),(3,1) })
(3,{1}/1;{ })



Output 1



Output 2

a

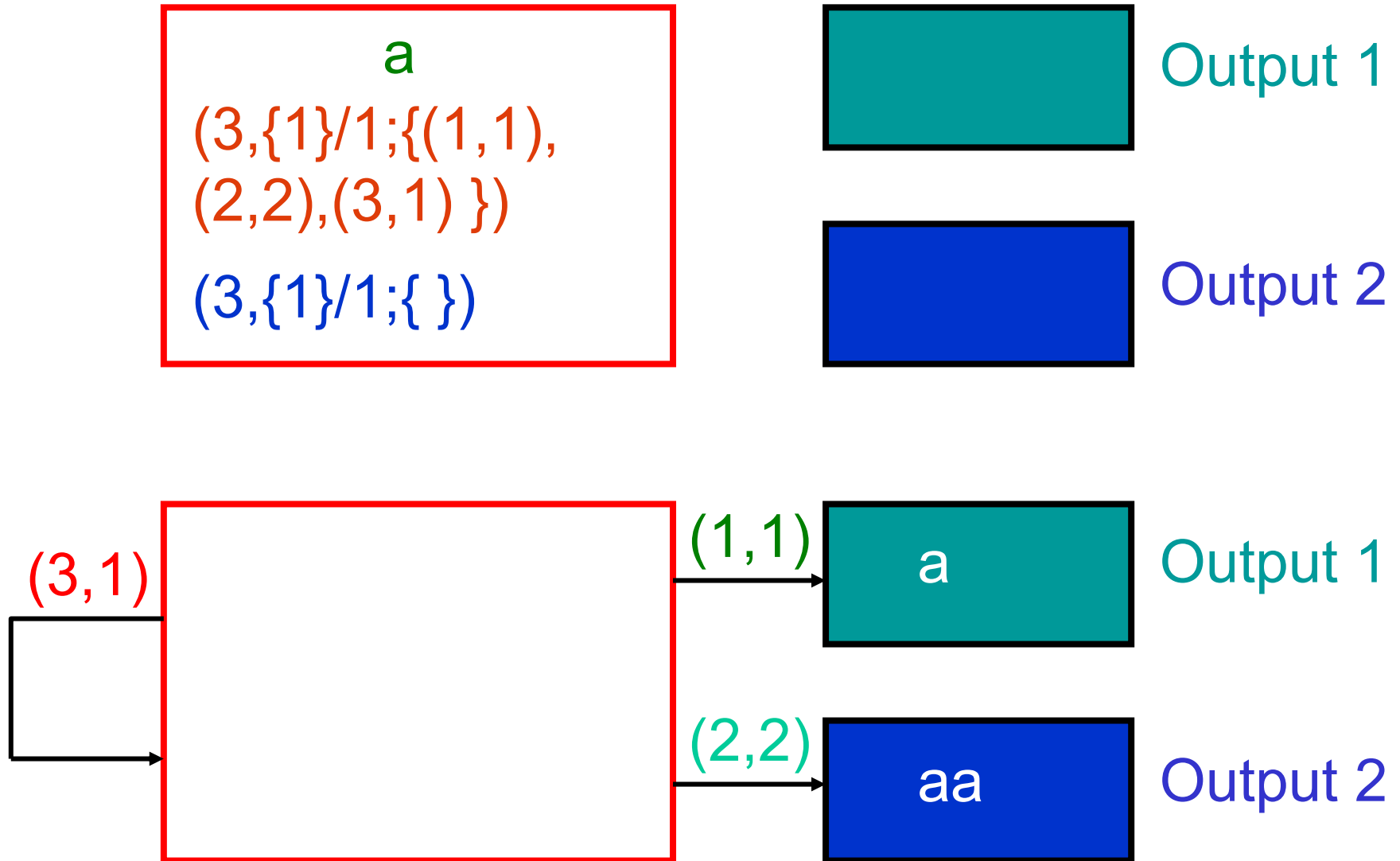


Output 1



Output 2

A First Example – Computation



A First Example – Computation

a
(3,{1}/1;{(1,1),
(2,2),(3,1) })
(3,{1}/1;{ })



Output 1



Output 2

a



Output 1



Output 2

A First Example – Computation

a
(3,{1}/1;{(1,1),
(2,2),(3,1) })
(3,{1}/1;{ })



Output 1



Output 2



aa

Output 1



aaaa

Output 2

Example : Characterization of Regular Sets of Non-negative Integers

A set of non-negative integers M is semilinear if and only if it can be generated by a finite (or bounded) ESNP system with only two neurons.

Let M be a semilinear set of non-negative integers generated by the regular grammar $G = (N, \{a\}, A_1, P)$, $N = \{A_i \mid 1 \leq i \leq m\}$, the start symbol A_1 , and P the set of regular productions of the forms $B \rightarrow aC$ and $A \rightarrow \lambda$ with $A, B, C \in N$.

We now construct the finite ESNP system $\Pi = (2, S, R)$ that generates an element of M by the number of spikes contained in the output neuron 1 at the end of a halting computation:

Example : Characterization of Regular Sets of Non-negative Integers

We start with one spike (representing the start symbol A_1 in neuron 2 (which keeps track of the actual non-terminal symbol) and no spike in the output neuron 1, i.e., $S = \{(1,0),(2,1)\}$. The production $A_i \rightarrow aA_j$ is simulated by the rule $(2, \{i\}/i; \{(1,1),(2,j)\})$ and $A_i \rightarrow \lambda$ is simulated by the rule $(2, \{i\}/i; \{ \})$, i.e., in sum we obtain

$$\Pi = (2, \{(1,0),(2,1)\}, R),$$

$$R = \{(2, \{i\}/i; \{(1,1),(2,j)\}) \mid 1 \leq i, j \leq m, A_i \rightarrow aA_j \in P\} \\ \cup \{(2, \{i\}/i; \{ \}) \mid 1 \leq i \leq m, A_i \rightarrow \lambda \in P\}.$$

Example : Characterization of Regular Sets of Non-negative Integers

We can also generate the numbers in M as the difference between the (first) two spikes arriving in the output neuron by the following ESNP system Π' :

$$\Pi' = (2, \{(1,0), (2, m+1)\}, R'),$$

$$R' = \{(2, \{m+1\}/m; \{(1,1), (2,1)\})\}$$

$$\cup \{(2, \{i\}/i; \{(2,j)\}) \mid 1 \leq i, j \leq m, A_i \rightarrow aA_j \in P\}$$

$$\cup \{(2, \{i\}/i; \{(1,1)\}) \mid 1 \leq i \leq m, A_i \rightarrow \lambda \in P\}.$$

Results: Bounded/Finite ESNP Systems

Lemma: For any ESNP system where during a computation only a bounded number of spikes occurs in the actor neurons, the generated language is regular.

The question whether in an ESNP system during a computation only a bounded number of spikes occurs in the actor neurons is not decidable, but:

Theorem: A language L with $L \subseteq T^*$ for a terminal alphabet T with $\text{card}(T)=n$ is regular if and only if it can be generated by a **finite** ESNP system with $n+1$ neurons.

Results: Bounded/Finite ESNP Systems

Corollary: A set of n -dimensional vectors is semilinear if and only if it can be generated by a finite ESNP system with $n+1$ neurons.

The preceding results depend on the way the output is taken – as the **number of spikes** present in the output neurons at the end of a halting computation and not as the **time between the first two spikes** in each output neuron.

Results: Unbounded ESNP Systems

Theorem: Any recursively enumerable language L with $L \subseteq T^*$ for a terminal alphabet T with $\text{card}(T)=n$ can be generated by an ESNP system with $n+3$ neurons.

Corollary: Any recursively enumerable set of n -dimensional vectors can be generated by an ESNP system with $n+3$ neurons.

Ideas for Future Research

- investigate **variants of definitions and/or restrictions** allowing for the characterization of families (of vectors of non-negative integers or strings) between the families of regular and recursively enumerable sets
- consider **inhibiting spikes**, i.e., allow negative values for the weights w in the productions (l, w, t) ; when this package arrives at the target cell l , then it causes this cell to be closed for t time steps (all such negative values are summed up at the moment they arrive at the cell provided it is not closed)

Thank You for Your Attention !